



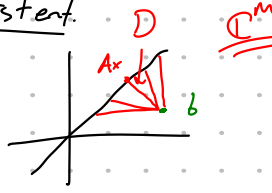
$$= \begin{pmatrix} 0 & 0 & 1 \\ 1 & 0 & 0 \\ 0 & 1 & 0 \end{pmatrix} \begin{pmatrix} \frac{1}{2} & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix} \begin{pmatrix} -0.8 & 0.6 & 0 & 0 \\ 0.6 & 0.8 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix}$$

$$= \begin{pmatrix} 0 & 0 & 1 \\ 1 & 0 & 0 \\ 0 & 1 & 0 \end{pmatrix} \begin{pmatrix} -0.4 & 0.3 & 0 & 0 \\ 0.6 & 0.8 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix} = \begin{pmatrix} 0 & 0 & 0 & 0 \\ -0.4 & 0.3 & 0 & 0 \\ 0.6 & 0.8 & 0 & 0 \end{pmatrix}$$

Inconsistent and Undetermined Systems (Say  $A$  is invertible, then  $Ax=b \Rightarrow x=A^{-1}b$ )

Consider  $Ax=b$  for an  $m \times n$  matrix  $A$ .  
Then, the minimal solution of this system is the least squares solution that has the least Euclidean norm if the system is inconsistent and it is the exact solution with the least Euclidean norm if the system is consistent.

$$A \in \mathbb{C}^{m \times n} \Rightarrow L_A: \mathbb{C}^n \rightarrow \mathbb{C}^m \quad L_A(x) = Ax$$



Alternative definition: Let  $D = \inf_{x \in \mathbb{C}^n} \|Ax - b\|_2$

Then the minimal solution of  $Ax=b$  is the (unique) element  $x_0 \in \mathbb{C}^n$  satisfying  $\|x_0\|_2 = \inf \{ \|x\|_2 \mid x \in \mathbb{C}^n, \|Ax - b\|_2 = D \}$

Theorem The minimal solution of  $Ax=b$  is given by  $x = A^+b$ .

Example Find the minimal solution of

$$\begin{aligned} 0x - 1.6y + 0.6z &= 5 \\ 0x + 1.2y + 0.8z &= 7 \\ 0x + 0y + 0z &= 3 \\ 0x + 0y + 0z &= -2 \end{aligned} \quad \rightarrow \quad \begin{matrix} A \\ \begin{pmatrix} 0 & -1.6 & 0.6 \\ 0 & 1.2 & 0.8 \\ 0 & 0 & 0 \\ 0 & 0 & 0 \end{pmatrix} \begin{pmatrix} x \\ y \\ z \end{pmatrix} = \begin{pmatrix} 5 \\ 7 \\ 3 \\ 2 \end{pmatrix} \end{matrix}$$

$$\Rightarrow \begin{pmatrix} x \\ y \\ z \end{pmatrix} = \begin{pmatrix} 0 & -1.6 & 0.6 \\ 0 & 1.2 & 0.8 \\ 0 & 0 & 0 \\ 0 & 0 & 0 \end{pmatrix}^+ \begin{pmatrix} 5 \\ 7 \\ 3 \\ 2 \end{pmatrix}$$

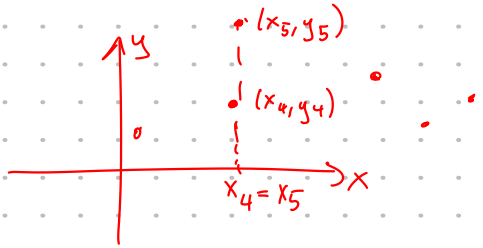
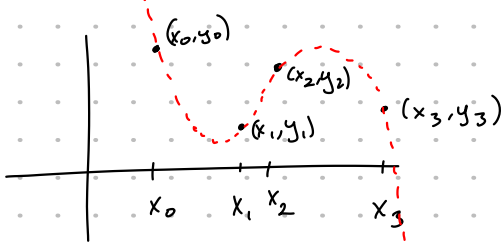
$$= \begin{pmatrix} 0 & 0 & 0 & 0 \\ -0.4 & 0.3 & 0 & 0 \\ 0.6 & 0.8 & 0 & 0 \end{pmatrix} \begin{pmatrix} 5 \\ 7 \\ 3 \\ 2 \end{pmatrix} = \begin{pmatrix} 0 \\ 0.1 \\ 8.6 \end{pmatrix}$$

## Chp 6 Approximating Functions

### 6.1 Polynomial Interpolation

Theorem If  $x_0, x_1, \dots, x_n$  are distinct real numbers, then for arbitrary values  $y_0, y_1, \dots, y_n$ ,  $\exists!$  polynomial  $P_n$  of degree at most  $n$  such that

$$P_n(x_i) = y_i \quad (0 \leq i \leq n)$$



Pf We prove uniqueness first.

Say  $p_n$  and  $q_n$  are two polynomials satisfying the conditions of the theorem. Then,

$$p_n(x_i) = y_i = q_n(x_i) \Rightarrow (p - q)(x_i) = 0 \quad 0 \leq i \leq n$$

Note that  $p - q$  is a polynomial of degree at most  $n$  but it has  $n+1$  zeros. Thus, it must be the constant 0 polynomial.

Next existence: We use induction on  $n$ .

For  $n=0$ , set  $p(x) = y_0$  ( $\deg p = 0$ )  $\checkmark$

Assume that the statement is true for  $n = k-1$ .

Given  $x_0, x_1, \dots, x_{k-1}, x_k, y_0, \dots, y_{k-1}, y_k$

by induction hypothesis,  $\exists p_{k-1}$  of degree  $\leq k-1$  s.t.

$$p_{k-1}(x_i) = y_i \quad (0 \leq i \leq k-1)$$

Define  $p_k(x) = p_{k-1}(x) + c \underbrace{(x-x_0)(x-x_1)\dots(x-x_{k-1})}_{(x_i = x_i)}$  for some  $c \in \mathbb{R}$  (to be determined below)

Clearly  $\deg(p_k) \leq k$ ,

$$p_k(x_i) = p_{k-1}(x_i) = y_i \quad \text{for } i \leq k-1.$$

$$\text{and } p_k(x_k) = p_{k-1}(x_k) + c(x_k - x_0)(x_k - x_1)\dots(x_k - x_{k-1}) = y_k \quad (*)$$

Thus, it is clear that  $\exists c \in \mathbb{R}$  s.t.  $p_k(x_k) = y_k$ .

Note that the proof gives us

$$p_0(x) = c_0$$

$$p_1(x) = c_0 + c_1(x - x_0)$$

$$p_2(x) = c_0 + c_1(x - x_0) + c_2(x - x_0)(x - x_1)$$

$\vdots$

$$p_n(x) = c_0 + c_1(x - x_0) + \dots + c_n(x - x_0)(x - x_1)\dots(x - x_{n-1})$$

(for some appropriate choices of  $c_i$ )

These polynomials are called the interpolation polynomials in Newton's form.

To evaluate  $P_n(x)$ , we will use Horner's algorithm.

Recall (Horner's Algorithm)  $d_0 = (x - x_0)$   $d_1 = (x - x_1) \dots$

If  $u = c_0 + c_1 d_0 + c_2 d_0 d_1 + \dots + c_k d_0 d_1 \dots d_{k-1}$

then  $u = (\dots ((c_k d_{k-1} + c_{k-1}) d_{k-2} + c_{k-2}) d_{k-3} + \dots + c_1) d_0 + c_0$ .

We can compute  $u$  by the following series of operations

$$u_k \leftarrow c_k$$

$$u_{k-1} \leftarrow u_k d_{k-1} + c_{k-1}$$

$$u_{k-2} \leftarrow u_{k-1} d_{k-2} + c_{k-2}$$

$\vdots$

$$u = u_0 \leftarrow u_1 d_0 + c_0$$

or more compactly by,

$$u \leftarrow c_k$$

for  $i = k-1$  to 0 step -1 do

$$u \leftarrow u d_i + c_i$$

end do

for our polynomial,  $d_i = (x - x_i)$   $[P_n(x) = c_0 + c_1(x - x_0) + \dots + c_n(x - x_0) \dots (x - x_{n-1})]$

Thus, we can compute  $u = P_k(t)$  (for a fixed  $t$ ) using

$$u \leftarrow c_k$$

for  $i = k-1$  to 0 step -1 do

$$u \leftarrow (t - x_i)u + c_i \quad d_i = (t - x_i)$$

end do

Note that  $c_k = \frac{y_k - P_{k-1}(x_k)}{(x_k - x_0)(x_k - x_1) \dots (x_k - x_{k-1})}$  using  $\otimes$

which we can compute by using similar algorithms but we will talk about a faster way to compute them later.

Example Find the lowest degree polynomial that passes through  $(2, 2)$ ,  $(3, 8)$ , and  $(5, 0)$ .

i.e.

$i$	0	1	2
$x_i$	2	3	5
$y_i$	2	8	0

$$P_0(x) = c_0 = y_0 = 2$$

$$P_1(x) = P_0(x) + c_1(x - x_0) = 2 + c_1(x - 2)$$

$$P_1(3) = 8 \Rightarrow 2 + c_1(3 - 2) = 8 \Rightarrow c_1 = 6$$

$$P_2(x) = 2 + 6(x - 2) + c_2(x - 2)(x - 3)$$

$$P_2(5) = 0 \Rightarrow 2 + 6(3) + c_2(3)(2) = 0$$

$$6c_2 = -20 \quad c_2 = -\frac{10}{3}$$

Thus,  $P(x) = P_2(x) = 2 + 6(x - 2) - \frac{10}{3}(x - 2)(x - 3)$

$$P(x) = \left( \left( -\frac{10}{3} \right) (x - 3) + 6 \right) (x - 2) + 2$$

$$P(2) = ( \quad ) (2 - 2) + 2 = 2$$

$$P(3) = \left( -\frac{10}{3}(3 - 3) + 6 \right) (3 - 2) + 2 = 8$$

$$P(5) = \left( -\frac{10}{3}(5 - 3) + 6 \right) (5 - 2) + 2 = \left( -\frac{20}{3} + 6 \right) 3 + 2 = 0$$