

Common mistakes so far (not all relevant to the midterm exam)

Order of quantifiers: (Example)

$f: (a,b) \rightarrow \mathbb{R}$ satisfies P1 if $\exists \lambda < 1$ s.t. $\forall x \in (a,b) f(x) \leq \lambda$
 \equiv P2 if $\forall x \in (a,b) \exists \lambda < 1$ s.t. $f(x) \leq \lambda$
 (implicitly depends on x)

equivalently, we can write

P2: if $\forall x \in \mathbb{R}, \exists \lambda_x < 1$ s.t. $f(x) \leq \lambda_x$ to show the dependence on x more explicitly.

a) Prove that if f satisfies P1 then it satisfies P2.

b) Find examples of f that satisfy P2 but not P1.

We want to show that f satisfies P2:

Given $x \in \mathbb{R}$ (we need to find $\lambda_x < 1$ s.t. $f(x) \leq \lambda_x$). Since f satisfies P1, if (this color is green) then $x^2 > -1$ for $x \in \mathbb{R}$

$\exists \lambda < 1$ s.t. $f(x) \leq \lambda \forall x \in \mathbb{R}$. All we have to do is

set $\lambda_x = \lambda$ thus $f(x) \leq \lambda = \lambda_x$ (a) ✓
 needed in P2 comes from P1

b) $\frac{1}{x}$ $y = \frac{1}{x}$ on $(1, \infty)$ $y = f(x)$

for any $x > 1$, $\frac{1}{x} < 1$. Set $\lambda_x = \frac{1}{x} < 1$ first found a suitable λ_x then proved P2

then $\forall x > 1, \lambda_x = \frac{1}{x}$ so $f(x) = \frac{1}{x} \leq \lambda_x$ P2 ✓

P1 is not satisfied since $\forall 0 < \lambda < 1, \frac{1}{\lambda} > 1 \Rightarrow \exists x$ s.t. $\frac{1}{\lambda} > x > 1$
 $x \in (1, \infty)$ but $f(x) = \frac{1}{x} > \lambda$ $\lambda < \frac{1}{x}$

Technicalities of ϵ - δ proofs

$\lim_{x \rightarrow a} f(x) = L$ if

$(\forall \epsilon > 0 \exists \delta > 0$ s.t. $\forall x \in \mathbb{R}$ (that satisfies) $|x-a| < \delta, |f(x)-L| < \epsilon$) is true

Note that δ may depend on ϵ but not on x .

More generally, The limit $\lim_{x \rightarrow a} f(x)$ exists if

$\exists L \in \mathbb{R}$ s.t. $(\forall \epsilon > 0 \exists \delta > 0$ s.t. $\forall x \in \mathbb{R}$ (that satisfies) $|x-a| < \delta, |f(x)-L| < \epsilon$) is true

We negate it as follows:

The limit $\lim_{x \rightarrow a} f(x)$ does not exist if

$\forall L \in \mathbb{R}$ $(\forall \epsilon > 0 \exists \delta > 0$ s.t. $\forall x \in \mathbb{R}$ (that satisfies) $|x-a| < \delta, |f(x)-L| < \epsilon$) is false

Example Prove that $\lim_{x \rightarrow 1} 2x+5 = 7$. if P1 then P2

We cannot (in general) instantly start proving the formal definition

$(\forall \epsilon > 0 \exists \delta > 0$ s.t. $\forall x \in \mathbb{R}$ (that satisfies) $|x-a| < \delta, |f(x)-L| < \epsilon$)

because we don't know what δ works. Instead we first do some work to figure out a suitable δ and then prove the formal statement using a δ that works!

We want

$|2x+5-7| < \epsilon$ whenever $0 < |x-1| < \delta$ reduce it to this

$|2x-2| = 2|x-1| < \epsilon$

so $|x-1| < \frac{\epsilon}{2}$. Thus, it seems like $\delta = \frac{\epsilon}{2}$ should work.

figure out what δ works

actual proof

Given $\epsilon > 0$, let $\delta = \frac{\epsilon}{2} > 0$. Then

$|x-1| < \delta = \frac{\epsilon}{2} \Rightarrow 2|x-1| < \epsilon$

$\Rightarrow |2x-2| < \epsilon$

$\Rightarrow |2x+5-7| < \epsilon$

i.e. $|f(x)-L| < \epsilon$ which was to be shown.

When computing $fl(x)$

$u = 2^{-23} \quad e = \frac{3}{4}u$
 $1+e = 1$ but $1-e \neq 1$

First, you need to find the binary representation of x .

e.g. $x = \frac{5}{32} = (0.00101)_2$ (fixed point)

Then write x in the form $x = (1.f)_2 \times 2^m$ (floating point)

e.g. $\frac{5}{32} = (1.01)_2 \times 2^{-3}$

Then we truncate the bits after the 23rd digit to the right of the binary point in floating point representation, to get x_- .

To get x_+ , we add 1 to the 23rd digit, that is,

$x_+ = x_- + 2^{-23} \cdot 2^m = x_- + 2^{-26}$ mach. numbers \downarrow

Then check if x is closer to x_+ or x_- to determine which one is $fl(x)$.

Chp 5 Selected Topics in Numerical Linear Algebra

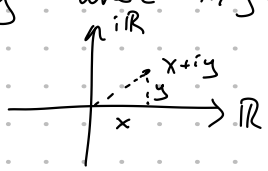
Review of Basic Concepts

\mathbb{C} : Complex numbers $z = x + iy$ where $x, y \in \mathbb{R}$ and $i = \sqrt{-1}$ that is $i^2 = -1$

conjugate: $\bar{z} = x - iy$

modulus: $|z| = \sqrt{x^2 + y^2}$

Notice that $\overline{\bar{z}} = z$ and $z\bar{z} = |z|^2$



$x^2 + 1$: no roots over \mathbb{R} .

Fundamental Theorem of Algebra

Over \mathbb{C} , every non-constant polynomial has at least one root.

If f is polynomial of degree n , then it is a product of n linear terms

$$f(z) = \lambda(z-z_1)(z-z_2)\dots(z-z_n)$$

The (standard) inner product and the Euclidean norm on \mathbb{C}^n

is defined by

$$\langle x, y \rangle = \sum_{j=1}^n x_j \bar{y}_j = x^T \bar{y}$$

$$\|x\|_2 = \sqrt{\langle x, x \rangle} \leftarrow \text{real valued quantity}$$

Notice that

$$\langle x, y \rangle = \overline{\langle y, x \rangle} \quad \text{and} \quad \langle x, \lambda y \rangle = \bar{\lambda} \langle x, y \rangle$$

\uparrow
 \mathbb{C} -anti-linear.

$$\text{and } \langle x+y, z \rangle = \langle x, z \rangle + \langle y, z \rangle$$

$\langle x, y \rangle$ is \mathbb{C} -linear in x and \mathbb{C} -anti-linear in y

If A is a complex $n \times k$ matrix then A^* is the (complex)

$k \times n$ matrix

$$A^* = \overline{A}^T \quad \text{that is } (A^*)_{jk} = \overline{A_{kj}}$$

Note that $\langle x, y \rangle = x^T \bar{y} = (x^T \bar{y})^T = \bar{y}^T x^T = y^* x$

and

$$x^* x = \langle x, x \rangle = \|x\|_2^2 = \sum_{j=1}^n x_j \bar{x}_j = \sum_{j=1}^n |x_j|^2$$